

Lecture 14.

This week we will switch to the fantasy genre: cloning, teleportation, etc. A pivotal role in these algorithms is played by EPR pairs or, more generally, entangled states.

Def-n. Let V and W be two finite dimensional vector spaces. A vector $h \in V \otimes W$ is called indecomposable provided h cannot be written as $h = v \otimes w$ with $v \in V$ and $w \in W$.

Similarly, let $|z\rangle \in \mathcal{H}_1 \otimes \mathcal{H}_2$ be a state vector. Then $|z\rangle$ is called entangled provided $|z\rangle$ cannot be written as $|z\rangle = |y\rangle \otimes |w\rangle$ with $|y\rangle \in \mathcal{H}_1$ and $|w\rangle \in \mathcal{H}_2$.

Example. Let $\mathcal{H}_1 = \mathcal{H}_2 = \mathbb{C}^2$ and $|z\rangle = \frac{1}{\sqrt{2}}(|00\rangle + |11\rangle)$. We will show that $|z\rangle$ is entangled. Indeed, otherwise

$$|z\rangle = \frac{1}{\sqrt{2}}(|00\rangle + |11\rangle) = (\alpha_0|0\rangle + \alpha_1|1\rangle) \otimes (\beta_0|0\rangle + \beta_1|1\rangle) = \alpha_0\beta_0|00\rangle + \alpha_0\beta_1|01\rangle + \alpha_1\beta_0|10\rangle + \alpha_1\beta_1|11\rangle \text{ meaning}$$

$$\begin{cases} \alpha_0\beta_0 = \frac{1}{\sqrt{2}} \\ \alpha_1\beta_0 = \alpha_0\beta_1 = 0 \\ \alpha_1\beta_1 = \frac{1}{\sqrt{2}} \end{cases}$$

It is easy to see that the system above has no solutions. Rmk. The vector $|z\rangle$ above is called an EPR pair (after Einstein, Podolsky and Rosen, see their '1935 paper' on Canvas) or Bell state (see his '1964 paper' on Canvas).

No-cloning theorem. There is no operator $U \in U(\mathbb{C}^2)^{\otimes n}$ such that $U(|\psi\rangle \otimes |0^{n-1}\rangle) = |\psi\rangle \otimes |\psi\rangle \otimes |\text{garbage}\rangle$ for all qubits $|\psi\rangle \in \mathbb{C}^2$.
↑ ancillas $(\mathbb{C}^2)^{\otimes n-2}$

Proof. Straightforward: suppose such an operator U exists,

then $U(|0\rangle \otimes |0^{n-1}\rangle) = |0\rangle \otimes |0\rangle \otimes |\text{garbage}_0\rangle$

$U(|1\rangle \otimes |0^{n-1}\rangle) = |1\rangle \otimes |1\rangle \otimes |\text{garbage}_1\rangle$

and, as U is linear,

$$U(|+\rangle \otimes |0^{n-1}\rangle) = \frac{1}{\sqrt{2}} (U(|0\rangle \otimes |0^{n-1}\rangle) + U(|1\rangle \otimes |0^{n-1}\rangle)) =$$

$$= \frac{1}{\sqrt{2}} (|0\rangle \otimes |0\rangle \otimes |\text{garbage}_0\rangle + |+\rangle \otimes |+\rangle \otimes |\text{garbage}_+\rangle).$$

On the other hand, we would like to have

$$U(|+\rangle \otimes |0^{n-1}\rangle) = |+\rangle \otimes |+\rangle \otimes |\text{garbage}_+\rangle.$$

Let's show that the equality

(*) $\frac{1}{\sqrt{2}} (|0\rangle \otimes |0\rangle \otimes |\text{garbage}_0\rangle + |1\rangle \otimes |1\rangle \otimes |\text{garbage}_1\rangle) = |+\rangle \otimes |+\rangle \otimes |\text{garbage}_+\rangle$
 can not hold true.

Indeed, $|+\rangle \otimes |+\rangle \otimes |\text{garbage}_+\rangle = \frac{1}{2} (|00\rangle + |01\rangle + |10\rangle + |11\rangle) \otimes |\text{garbage}_+\rangle$

Measuring the first two qubits on the l.h.s. and r.h.s. of (*), we see that probabilities do not match up. For instance, the probability of getting $|00\rangle$ on the l.h.s. is 0, while it is $(\frac{1}{2})^2 = \frac{1}{4}$ on the r.h.s.

Kmk. We have actually shown that it is impossible to 'clone' the three state vectors $|0\rangle$, $|1\rangle$ and $|+\rangle$ simultaneously.

Quantum teleportation.

Suppose Alice has a qubit $|\psi\rangle = \alpha|0\rangle + \beta|1\rangle$, which she would like Bob to have. Let's assume that Alice and Bob have prepared an EPR pair $\frac{1}{\sqrt{2}}(|00\rangle + |11\rangle)$, Alice stored the first qubit and Bob stored the second. Moreover, Alice can text Bob two classical bits.

Here is a way to 'teleport' $|\psi\rangle$.

Step 1. Alice has two qubits: $|\psi\rangle$ and the first qubit of the EPR pair. She applies CNOT with $|\psi\rangle$ being the control bit:

$$|\psi\rangle \otimes \frac{1}{\sqrt{2}}(|00\rangle + |11\rangle) = (\alpha|0\rangle + \beta|1\rangle) \otimes \frac{1}{\sqrt{2}}(|00\rangle + |11\rangle) \rightsquigarrow$$

$$\rightsquigarrow \frac{1}{\sqrt{2}} (\alpha|000\rangle + \alpha|011\rangle + \beta|110\rangle + \beta|101\rangle).$$

Step 2. Alice measures her second qubit:

$$P(\text{Alice's second qubit is } |0\rangle) = \left(\frac{1}{\sqrt{2}}\right)^2 (|\alpha|^2 + |\beta|^2) = \frac{1}{2}$$

↑
↑
from $|000\rangle$
from $|101\rangle$

$$P(\text{Alice's second qubit is } |1\rangle) = \left(\frac{1}{\sqrt{2}}\right)^2 (|\alpha|^2 + |\beta|^2) = \frac{1}{2} \text{ (or simply } 1 - \frac{1}{2} = \frac{1}{2}\text{)}$$

↑
↑
from $|011\rangle$
from $|110\rangle$

In the first case (2nd qubit is $|0\rangle$) the state collapses to $\frac{1}{\sqrt{2}}(\alpha|00\rangle + \beta|11\rangle)$, while if the second qubit is $|1\rangle$, the

state collapses to $\frac{1}{\sqrt{2}}(\alpha|10\rangle + \beta|11\rangle)$, in which case Alice sends Bob the classical bit '1' and he applies a NOT to his state. This way, the resulting state at this point will always be $\frac{1}{\sqrt{2}}(\alpha|100\rangle + \beta|111\rangle)$.

Step 3. Alice applies the Hadamard operator to her qubit: $\frac{1}{\sqrt{2}}(\alpha|100\rangle + \beta|111\rangle) \rightarrow \frac{1}{\sqrt{2}}(\alpha|1+0\rangle + \beta|1-1\rangle) =$
 $= \frac{1}{\sqrt{2}}(\alpha|100\rangle + \alpha|110\rangle + \beta|101\rangle - \beta|111\rangle)$.

Step 4. Alice measures her qubit:

$$P(\text{Alice's qubit collapses to } |0\rangle) = \left(\frac{1}{\sqrt{2}}\right)^2 (\underbrace{|\alpha|^2}_{\text{from } |100\rangle} + \underbrace{|\beta|^2}_{\text{from } |101\rangle}) = \frac{1}{2}$$

$$P(\text{Alice's qubit collapses to } |1\rangle) = \frac{1}{2}$$

1st case: Bob has $\alpha|10\rangle + \beta|11\rangle = |\Psi\rangle \checkmark$

2nd case: Bob has $\alpha|10\rangle - \beta|11\rangle$. Then Alice sends Bob a 1-bit message '1' (2nd classical bit) and he applies the operator $Z = \begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix} \in U_2(\mathbb{C})$ on the qubit to get $Z(\alpha|10\rangle - \beta|11\rangle) = \alpha|10\rangle + \beta|11\rangle = |\Psi\rangle \checkmark$

Informal slogan: '1 ebit + 2 bits \geq 1 qubit' (Bennet's law)

↑
EPR pair

Rmk. Alice doesn't have $|\Psi\rangle$ anymore, so no-cloning thm was not violated.

Entanglement Swapping.

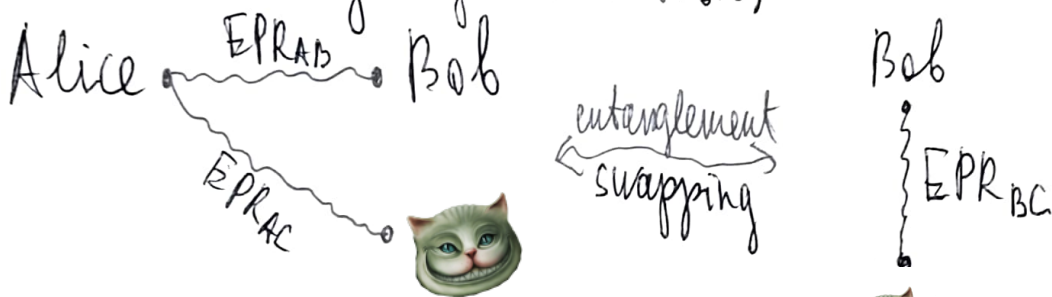
Suppose Alice shares an EPR pair with Bob

$$EPR_{AB} = \frac{1}{\sqrt{2}} (|00\rangle + |11\rangle)$$

and another EPR pair with Cheshire cat

$$EPR_{AC} = \frac{1}{\sqrt{2}} (|00\rangle + |11\rangle).$$

The entanglement swapping allows to create an entangled pair of qubits between Bob and the Cheshire cat (without them having any interaction):



The procedure is essentially build on Alice's teleportation to Bob her part (qubit) of the EPR_{AC} pair. We give an outline.

① Starting state: $\frac{1}{2} (|00\rangle + |11\rangle) \otimes (|00\rangle + |11\rangle) = \frac{1}{2} (|0000\rangle + |0011\rangle + |1000\rangle + |1011\rangle) + \frac{1}{2} (|1100\rangle + |1111\rangle)$ (the 1st and 3rd qubit belong to Alice).

② Alice applies $CNOT_{13}$ (1st qubit is controlling) giving the state $\frac{1}{2} (|0000\rangle + |0011\rangle + |1100\rangle + |1101\rangle)$.

③ Alice measures the third qubit producing the state $\frac{1}{\sqrt{2}} (|000\rangle + |1111\rangle) \rightarrow$ texts Bob the classical bit 'False' (0).

$\frac{1}{\sqrt{2}} (|100\rangle + |110\rangle) \rightarrow$ texts Bob a 'True' (1) to apply NOT (and 🐱)

As a result the state will be $\frac{1}{\sqrt{2}}(|000\rangle + |111\rangle)$ in both cases.

(4) Alice applies H on the 1st qubit, the state becomes

$$\frac{1}{\sqrt{2}}(|+00\rangle + |-11\rangle) = \frac{1}{2}(|000\rangle + |100\rangle + |011\rangle - |111\rangle)$$

(5) Alice measures her first qubit:

if $|0\rangle$, the state collapses to $\frac{1}{\sqrt{2}}(|00\rangle + |11\rangle)$;

if $|1\rangle$, the state collapses to $\frac{1}{\sqrt{2}}(|00\rangle - |11\rangle)$.

In the first case Alice texts both Bob and the cat 'False'; otherwise she texts one of them 'True' and the other 'False'. The one that receives 'True' applies Z to his qubit.